# Using Java-Prolog connection to implement automated reasoning \*

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Abstract. The concept of knowledge base with output was introduced in [12]. In this paper we present a synthesis of some results obtained by implementing a particular case, where the output space is a set of sentences in a natural language. The corresponding method uses a labelled graph and therefore a path driven reasoning is obtained. The inference engine is realized in Prolog and the graphical user interface is realized in Java. In order to realize this application the connection Java-Prolog written by Ugo Chirico is used.

**AMS Classification:** Primary: 68T30, Secondary: 68N17, 68N19 **Keywords:** inference engine, labelled graph, labelled stratified graph, Peano algebra, knowledge base with output, Java-Prolog connection

## 1 Introduction

Knowledge representation is a central problem and a major area of research in artificial intelligence. The most known methods for knowledge representation can be classified into two main categories. The first one uses the logical representation and the second category uses the graph- based representation of knowledge. There have been a great number of proposals for the semantics of logic representation of knowledge ([1], [4], [5], [17]). In a graph-based representation an entity is given by a pair of nodes and a link between them. The semantics is given by specifying the concrete meaning of this notational convention. In this category can be introduced the following structures: semantic networks, frames, scripts, conceptual graphs. An overview of these methods can be found in ([18]). Many of these methods, both from the first category and the second category, are implemented today ([2], [6], [8], [10], [11]).

The method implemented in this paper has established the starting point in my research work concerning the concept of knowledge base with output ([12]). Several intuitive aspects concerning both the syntactic and semantic computations in a knowledge base with output are described in [14]. We relieve that

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in the present paper the output space is a set of sentences in a natural language and therefore the corresponding application can be used in the domain of communicating semantics.

## 2 Theoretical results

In what follows we shall use the concepts of labelled graph and labelled stratified graph ([12], [13]).

We consider two finite sets S and  $L_0$  such that  $S \cap L_0 = \emptyset$ . An element of S is called a **node label**; the elements of  $L_0$  are called **arc labels**.

A binary relation over S is a subset  $\rho \subseteq S \times S$ . If  $\rho_1 \in 2^{S \times S}$  and  $\rho_2 \in 2^{S \times S}$ then we define:  $\rho_1 \circ \rho_2 = \{(x, y) \in S \times S \mid \exists z \in S : (x, z) \in \rho_1, (z, y) \in \rho_2\}$ . We define the mapping  $prod : dom(prod) \longrightarrow 2^{S \times S}$  as follows:

$$dom(prod) = \{(\rho_1, \rho_2) \in 2^{S \times S} \times 2^{S \times S} \mid \rho_1 \circ \rho_2 \neq \emptyset\}$$
  
$$prod(\rho_1, \rho_2) = \rho_1 \circ \rho_2$$

The mapping prod is called the *product operation*. The pair  $(2^{S \times S}, \sigma_S)$ , where  $\sigma_S = prod$  becomes a partial  $\sigma$ - algebra.

Let  $T_0$  be a set of binary relations on S, such that  $\emptyset \notin T_0$ . Let  $f_0 : L_0 \longrightarrow T_0$  be a surjective function. The system  $G = (S, L_0, T_0, f_0)$  is called a **labelled** graph.

A labelled graph  $G = (S, L_0, T_0, f_0)$  is represented as follows. Every node is represented by a rectangle containing its label. We draw a labelled arc  $a \in L_0$ from the node  $x \in S$  to the node  $y \in S$  if and only if  $(x, y) \in f_0(a)$  (Figure 1). Every element of S will designate a node of the graph and only one; thus some bijection between S and the set of the nodes can be established. We say that  $a \in L_0$  is a label of the binary relation  $f_0(a) \in T_0$ . Because  $f_0$  is a surjective function, it follows that every element of  $T_0$  has at least one label.



Fig. 1. A labelled arc

Let  $L_0$  be a nonempty set. Let  $\mathcal{H} = (H, \sigma_H)$  be the Peano  $\sigma$ -algebra over  $L_0$ . For every  $L \in Initial(\mathcal{H})$  such that  $L \supseteq L_0$  we define the partial  $\sigma$ -algebra  $\mathcal{A}_H(L) = (L, \sigma_L)$ , where:

- $dom(\sigma_L) = \{(x, y) \in L \times L \mid \sigma_H(x, y) \in L\}$
- $\sigma_L(x, y) = \sigma_H(x, y)$  for every  $(x, y) \in dom(\sigma_L)$

The pair  $\mathcal{A}_H(L) = (L, \sigma_L)$  is named the partial  $\sigma$ -algebra associated to  $L \in Initial(\mathcal{H})$ .

Let  $G = (S, L_0, T_0, f_0)$  be a labelled graph. A labelled stratified graph over G is a system  $\mathcal{G} = (G, L, T, \sigma_T, f)$  such that:

- $(T, \sigma_T) \in Env(T_0)$
- $L \in Initial(\mathcal{H})$  and  $L_0 \subseteq L$ , where  $\mathcal{H}$  is the Peano  $\sigma$ -algebra over  $L_0$
- $f : \mathcal{A}_H(L) \longrightarrow \mathcal{A}_T$  is a surjective morphism such that  $f_0 \prec f$  and if  $(f(u), f(v)) \in dom(\sigma_T)$  then  $\sigma_H(u, v) \in L$

Because  $\sigma$  is a symbol of arity 2, we have  $dom(\sigma_L) \subseteq L \times L$  and  $dom(\sigma_T) \subseteq T \times T$ . The mapping  $f : L \longrightarrow T$  is a morphism, that is, in the diagram



for every  $(u, v) \in dom(\sigma_L)$  the following two properties are satisfied:

 $(f(u), f(v)) \in dom(\sigma_T)$  $\sigma_T(f(u), f(v)) = f(\sigma_L(u, v))$ 

Let us consider a knowledge piece KP. In order to perform an automated reasoning based on KP, the following steps will be realized:

- the knowledge piece KP is transposed in a labelled graph G and its component  $S, L_0, T_0$  and  $f_0$  are obtained;
- based on [13] some labelled stratified graph  $(G, L, T, \sigma_T, f)$  is chosen;
- some output space O of sentences in a natural language is considered and the semantical computations described in [12] are realized

In order to illustrate the computations we shall consider the following knowledge piece KP: Peter is a student. John is a schoolboy. Every schoolboy is a presumable student. Peter is John's friend. Peter is George's brother. George is a student.

Taking the set of nodes  $S = \{John, schoolboy, Peter, student, George\}$  and the set of the initial labels  $L_0 = \{p, is\_a, pres, is\_brother\}$ , the labelled graph  $G = (S, L_0, T_0, f_0)$  represented in figure 2 is obtained, where

•  $T_0 = \{\rho_1, \rho_2, \rho_3, \rho_4\};$ 



Fig. 2. Labelled graph for KP

- $\rho_1 = \{(Peter, John)\}, \rho_2 = \{(John, schoolboy), (Peter, student), (George, student),$ student)},  $\rho_3 = \{(Peter, George)\}, \rho_4 = \{(schoolboy, student)\}$ •  $f_0(p) = \rho_1, f_0(is_a) = \rho_2, f_0(is_brother) = \rho_3, f_0(pres) = \rho_4$

The closure of  $T_0$  under  $prod_S$  is the set  $T = T_0 \cup \{\rho_5, \rho_6, \rho_7\}$ , where  $\rho_5 =$  $\{(Peter, schoolboy)\}, \rho_6 = \{(John, student)\} \text{ and } \rho_7 = \{(Peter, student)\}.$  Really, we have:

- $prod_S(\rho_1, \rho_2) = \rho_5, prod_S(\rho_2, \rho_4) = \rho_6$
- $prod_S(\rho_3, \rho_2) = \rho_7, prod_S(\rho_5, \rho_4) = \rho_7, prod_S(\rho_1, \rho_6) = \rho_7$

Taking into consideration the theoretical results exposed in the first part of this section we obtain the following computations:

$$\begin{aligned} f(\sigma(p, is_{-}a)) &= \rho_5, \ f(\sigma(is_{-}a, pres)) = \rho_6\\ f(\sigma(is_{-}brother, is_{-}a)) &= \rho_7, \ f(\sigma(p, \sigma(is_{-}a, pres))) = \rho_7, \\ f(\sigma(\sigma(p, is_{-}a), pres)) &= \rho_7 \end{aligned}$$

Based on these computations we can take for L the set obtained by adding to  $L_0$  the following elements:

 $\{\sigma(p, is\_a), \sigma(is\_a, pres), \sigma(is\_brother, is\_a), \sigma(p, \sigma(is\_a, pres)), \sigma(a, pres)\}$ 

On the other hand we observe that

$$f(\sigma(p, \sigma(is\_a, pres))) = f(\sigma(\sigma(p, is\_a), pres)) = \rho_7$$

and the semantics of the label  $\sigma(p, \sigma(is\_a, pres))$  is the same as the semantics of  $\sigma(\sigma(p, is\_a), pres)$ .

In conclusion, to obtain an environment for a labelled stratified graph (see [13]) we can choose the following mapping  $\sigma$ :

 $\sigma(p, is\_a) = a, \sigma(a, pres) = b, \sigma(is\_brother, is\_a) = c$  $\sigma(is\_a, pres) = d$ 

In the next section we present the manner in which the above results can be used to implement the corresponding reasoning.

## 3 Implementation

In order to implement a software product based on the results presented in the previous section, the **JIProlog** product written by Ugo Chirico was used (http://www.ugosweb.com/jiprolog). The name JIProlog derives from JavaInternetProlog and this name relieves the fact that a bidirectional connection Java-Prolog is realized. JIProlog is written in Java and the Prolog engine is compatible with other Prolog engines such as SWI Prolog. The user can invoke the Prolog interpreter in any Java applet/application and all the solutions of the Prolog engine can be obtained. A Prolog program can invoke in turn a Java method.

The implementation is based on the following components:

- the inference engine *rsem.pl* written in Prolog
- the component  $kp\_name.pl$
- the component *Transform.java*
- a graphical user interface ExtSem.java

The inference engine rsem.pl is able to realize the following:

- find the paths between two nodes of a labelled graph
- for each path computes the associated label
- gives a general structure of the explanation module

The component  $kp_name.pl$  realizes the following tasks:

- defines the labelled graph associated to a given knowledge piece
- receives the environment, that is the mapping  $\sigma$
- defines the semantics of the labels corresponding to the knowledge piece taken into consideration

Because a person is a node of the labelled graph and in Prolog such a name begins with a small letter, the component *Transform.java* receives a person name and transforms the first character in a capital letter. This component, written in Java, is invoked from Prolog to display the conclusion of the reasoning.

The graphical user interface uses the bidirectional connection Java-Prolog and realizes the following tasks:

- ten buttons to perform the following tasks:
  - $-\,$  a button to load the inference engine
  - $-\,$  a button to select some knowledge piece
  - a button to obtain the first conclusion and another button to obtain more conclusions of the reasoning
  - a button to reset the reasoning process
  - a button for help, buttons to display the text of the knowledge piece, the image of the labelled graph and the explanations for the conclusion
- five windows are used to display the conclusion of the reasoning and the explanations, the image of the labelled graph corresponding to the knowledge piece selected by button, the order of the actions, messages and the text of the knowledge piece



Fig. 3. The initial image of GUI



Fig. 4. Selecting a knowledge piece

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Fig. 5. First conclusion



Fig. 6. Conclusions and explanations

If we launch the graphical user interface (GUI), on the computer display the image from figure 3 is displayed.

The next step is the loading of the inference engine by acting the corresponding button. Then a knowledge piece must be selected and the image from figure 4 is shown on the display. The corresponding program in Prolog must be consulted by acting the item *Consult* of a menu. Using the button *LIST OF NODES* we can view all the nodes of the graph and in the next step we select both the first and the second node. By acting the button *FIRST CONCLUSION*, the image from figure 5 is exposed. Then by acting the button *EXPLANATIONS* and then *MORE CONCLUSIONS* the image from figure 6 is lay out.

A complete result of the reasoning corresponding to the first node *peter* and the second node *student* shows as follows:

CONCLUSION 1 Peter is the friend of a presumable student EXPLANATIONS Peter is the friend of John John is a schoolboy Every schoolboy is a presumable student CONCLUSION 2 Peter is the brother of a student EXPLANATIONS Peter is the brother of George George is a student CONCLUSION 3 Peter is a student EXPLANATIONS Peter is a student CONCLUSION 4 Finish solutions

## 4 Conclusions and future work

The idea used in this application has constituted the starting point to introduce and formalize the concept of knowledge base with output ([13]). As it is known, such a structure means a labelled stratified graph and an output space endowed with and algebraic operation. On the other hand, a labelled stratified graph uses a partial Peano algebra. Some algebraic morphism with values in the output space gives the semantics of the conclusion of some reasoning. In comparison with a semantic network, the labels are here abstract elements and these elements are "interpreted" in an output space. Various elements may be taken as elements of the output space. For example, the elements may be sentences in a natural language as we proceeded in this paper. The last research was focused on the algebraic properties of the labelled stratified graphs, especially to study so named *the choice problem* ([15], [16]). This problem means the choice of an appropriate labelled stratified graph for a given knowledge piece. The problem can be avoided by considering the greatest labelled stratified graph for a given context and then by introducing several constraints on the syntactic computations.

The future research in the implementation of these concepts includes the following:

1) consider the labelled stratified graph for a given context and formalize the constraints in the syntactic computations ([15], [16])

consider the case of an output space constituted from pieces of images and study the implications of this choice in image synthesis

consider the case of a non-associative binary operation from the output space  $% \left( {{{\bf{n}}_{\rm{a}}}} \right)$ 

We intend to study the applications of the concepts presented in this paper to problem solving, image synthesis and image analysis.

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